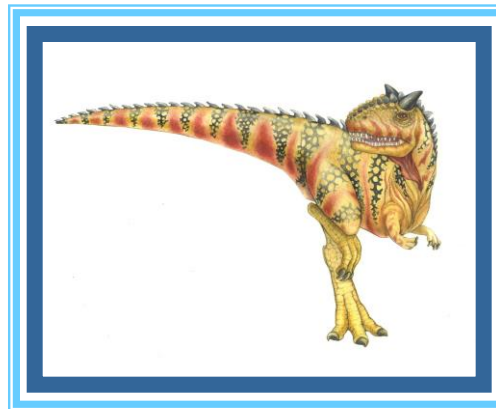
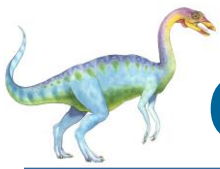


# Chapter 1: Introduction

---





# Chapter 1: Introduction

---

1. What Operating Systems Do
2. Computer-System Organization
3. Computer-System Architecture
4. Operating-System Structure
5. Operating-System Operations
6. Process Management
7. Memory Management
8. Storage Management
9. Protection and Security
10. Kernel Data Structures
11. Computing Environments
12. Open-Source Operating Systems

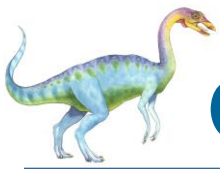




---

# 1.5 OPERATING-SYSTEM OPERATIONS



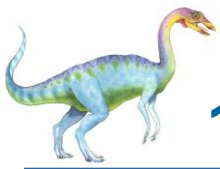


# Operating-System Operations

---

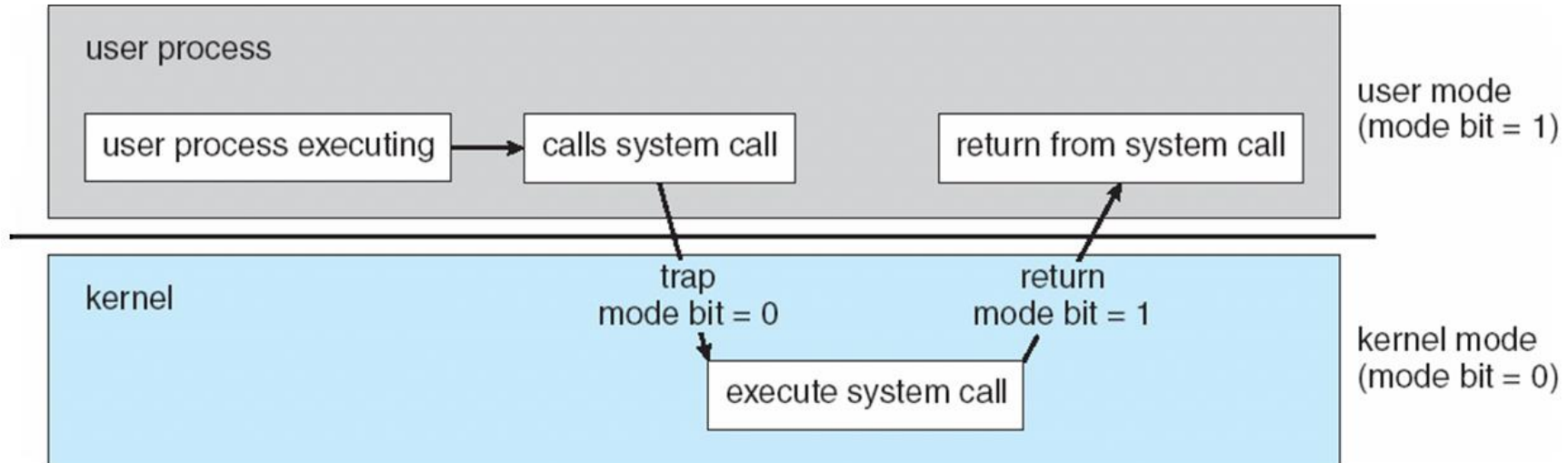
- **Interrupt driven** (hardware and software)
  - Hardware interrupt by one of the devices
  - Software interrupt (**exception** or **trap**):
    - ▶ Software error (e.g., division by zero)
    - ▶ Request for operating system service
    - ▶ Other process problems include infinite loop, processes modifying each other or the operating system

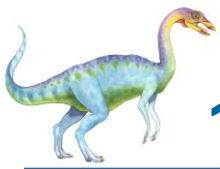




## 1.5.1 Dual-Mode and Multimode Operation (1)

- **Dual-mode** operation allows OS to protect itself and other system components
  - **User mode** and **kernel mode**(supervisor mode, system mode, privileged mode)



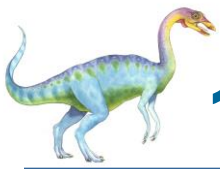


## 1.5.1 Dual-Mode and Multimode Operation (2)

---

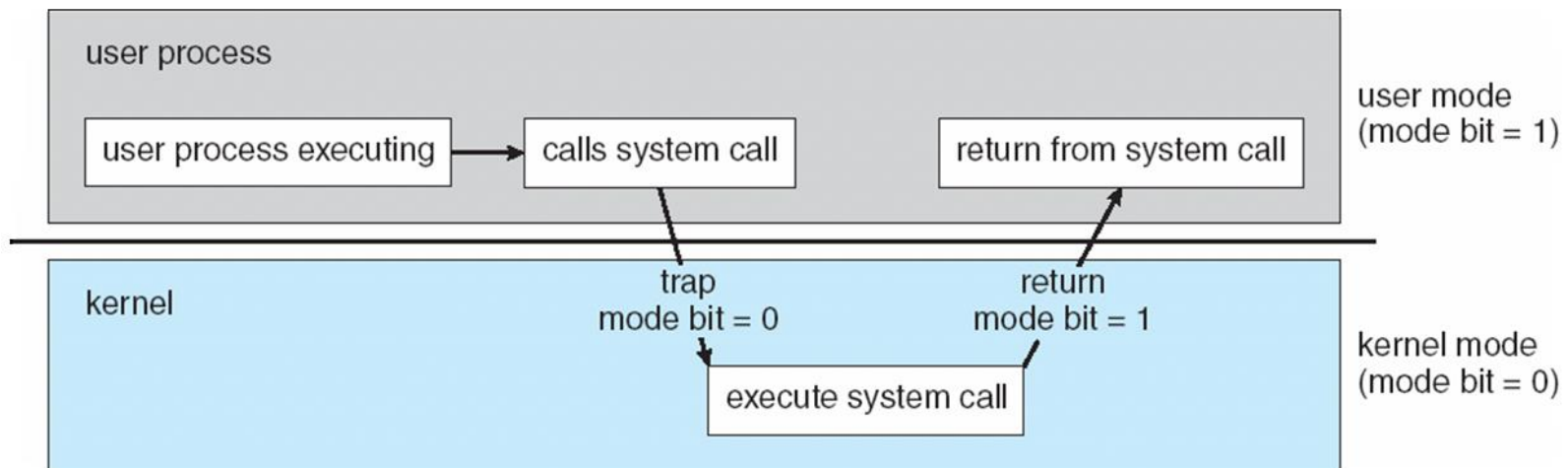
- Dual-mode (cont.)
  - **Mode bit** provided by hardware
    - ▶ Provides ability to distinguish when system is running user code or kernel code
    - ▶ Some instructions designated as privileged, only executable in kernel mode
    - ▶ System call changes mode to kernel, return from call resets it to user
  
- Increasingly CPUs support multi-mode operations
  - i.e. virtual machine manager (**VMM**) mode for guest VMs
  - User mode < VMM < kernel mode
  
- lack of a hardware-supported dual mode
  - cause serious shortcomings in an operating system





## 1.5.2 Timer

- Timer to prevent infinite loop / process hogging resources
  - Timer is set to interrupt the computer after some time period
  - Keep a counter that is decremented by the physical clock.
  - Operating system set the counter (privileged instruction)
  - When counter zero generate an interrupt
  - Set up before scheduling process to regain control or terminate program that exceeds allotted time





---

# 1.6 PROCESS MANAGEMENT





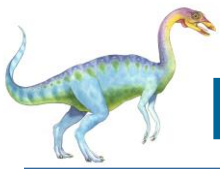


# Process Management

---

- A process is a program in execution. It is a unit of work within the system. Program is a **passive entity**, process is an **active entity**.
  
- Process needs resources to accomplish its task
  - CPU, memory, I/O, files
  - Initialization data
  
- Process termination requires reclaim of any reusable resources
  
- Single-threaded process has one **program counter** specifying location of next instruction to execute
  - Process executes instructions sequentially, one at a time, until completion
  
- Multi-threaded process has one program counter per thread
  
- Typically system has many processes, some user, some operating system running concurrently on one or more CPUs
  - Concurrency by multiplexing the CPUs among the processes / threads





# Process Management Activities

---

- The operating system is responsible for the following activities in connection with process management:
  - **Creating** and **deleting** both user and system processes
  - **Suspending** and **resuming** processes
  - Providing mechanisms for process **synchronization**
  - Providing mechanisms for process **communication**
  - Providing mechanisms for **deadlock** handling

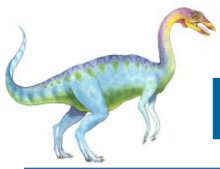




---

# 1.7 MEMORY MANAGEMENT





# Memory Management

---

- To execute a program **all (or part) of the instructions** must be in memory
- **All (or part) of the data** that is needed by the program must be in memory.
- Memory management determines **what is in memory and when**
  - Optimizing CPU utilization and computer response to users
- Memory management activities
  - **Keeping track** of which parts of memory are currently being used and by whom
  - **Deciding** which processes (or parts thereof) and data to move into and out of memory
  - **Allocating** and **deallocating** memory space as needed

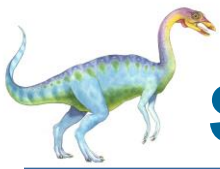




---

# 1.8 STORAGE MANAGEMENT





# Storage Management

---

- OS provides uniform, logical view of information storage
  - Abstracts physical properties to logical storage unit - **file**
  - Each medium is controlled by device (i.e., disk drive, tape drive)
    - ▶ Varying properties include access speed, capacity, data-transfer rate, access method (sequential or random)



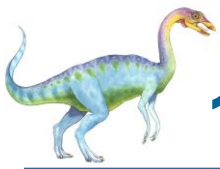


## 1.8.1 File-System management

---

- Files usually organized into **directories**
- **Access control** on most systems to determine who can access what
- OS activities include
  - **Creating** and **deleting** files and directories
  - Primitives to **manipulate** files and directories
  - **Mapping** files onto secondary storage
  - **Backup** files onto stable (non-volatile) storage media





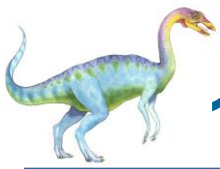
## 1.8.2 Mass-Storage Management

---

- Usually disks used to store data that does not fit in main memory or data that must be kept for a “long” period of time
- Proper management is of central importance
- Entire speed of computer operation **hinges on disk subsystem** and its algorithms
- OS activities
  - **Free-space management**
  - **Storage allocation**
  - **Disk scheduling**
- Some storage need not be fast
  - Tertiary storage includes optical storage, magnetic tape
  - Still must be managed – by OS or applications
  - Varies between WORM (write-once, read-many-times) and RW (read-write)





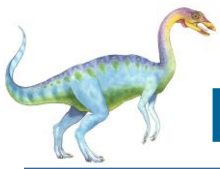


## 1.8.3 Caching

---

- Information is normally kept in some storage system (such as main memory)
  - As it is used, it is copied into a **faster** storage system—the cache—on a temporary basis.
- **internal programmable registers**, such as index registers, provide a high-speed cache for main memory
- Other caches are implemented totally in hardware.
  - Most systems have an **instruction cache** to hold the instructions expected to be executed next.
- Because caches have limited size, **cache management** is an important design problem

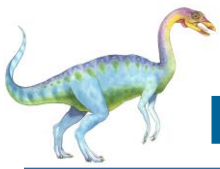




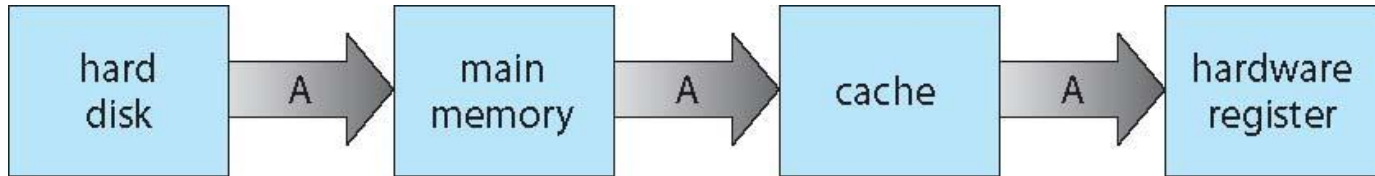
# Performance of Various Levels of Storage

Level	1	2	3	4	5
Name	registers	cache	main memory	solid state disk	magnetic disk
Typical size	< 1 KB	< 16MB	< 64GB	< 1 TB	< 10 TB
Implementation technology	custom memory with multiple ports CMOS	on-chip or off-chip CMOS SRAM	CMOS SRAM	flash memory	magnetic disk
Access time (ns)	0.25 - 0.5	0.5 - 25	80 - 250	25,000 - 50,000	5,000,000
Bandwidth (MB/sec)	20,000 - 100,000	5,000 - 10,000	1,000 - 5,000	500	20 - 150
Managed by	compiler	hardware	operating system	operating system	operating system
Backed by	cache	main memory	disk	disk	disk or tape



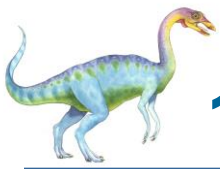


# Migration of integer A from disk to register



- The movement of information between levels of a storage hierarchy may be either explicit or implicit, depending on the hardware design and the controlling operating-system software
- Multitasking environments must be careful to use most recent value, no matter where it is stored in the storage hierarchy
- Multiprocessor environment must provide **cache coherency** in hardware such that all CPUs have the most recent value in their cache
- Distributed environment situation even more complex
  - Several copies of a datum can exist





## 1.8.4 I/O Systems

---

- One purpose of OS is to hide peculiarities of hardware devices from the user
  
- I/O subsystem responsible for
  - **Memory management of I/O** including
    - ▶ buffering : storing data temporarily while it is being transferred,
    - ▶ caching :storing parts of data in faster storage for performance,
    - ▶ spooling :the overlapping of output of one job with input of other jobs
  - **General device-driver interface**
  - Drivers for **specific** hardware devices

